# Theory of Computation Countability and Diagonalization

Arjun Chandrasekhar

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- We will show that there are strictly more languages than there are Turing machines
- ► This will imply that there is not a Turing machine for every language

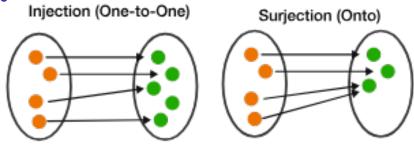
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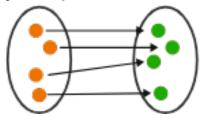
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- **Surjective:** Every element from  $S_2$  is mapped to at least once
- ▶ **Injective:** Every element of  $S_1$  maps to <u>exactly</u> one element of  $S_2$



#### Bijection (One-to-One and Onto)



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- ▶ A set *S* is **countably infinite** if there exists a bijection  $\mathbb{N} \mapsto S$
- Can also think of it as follows: can we write a program to print out the elements of S one by one, such that every element eventually gets printed if we let the program run long enough?

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  - Given enough time, every square number will (eventually) be printed

Let's prove that the set of integers  $\mathbb{Z}=\{\ldots,-3,-2,-1,0,1,2,3,\ldots\}$  is countably infinite

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Let's prove that the following set is countably infinite

$$\mathbb{N}^2 = \{(x, y) | x, y \in \mathbb{N}\}\$$

i.e. every combination of 2 natural numbers

**Hint:** for every integer k, there are only a finite number of (x, y) pairs that add up to k

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- ► Go through all combinations that add up to 1 ►  $1 \mapsto (1,0)$

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  - $\triangleright$  5  $\mapsto$  (0, 2)

# Countability of $\mathbb{N}^2$ : dovetailing

New "set" of combinations that add up to a common sum total

	0	1	2	3			
0	(0, 0)		(0, 2)	(0,3)			
1	(1, 0)	(1, 2)	(1, 2)	(1, 3)			
2	(2, 0)	(2, 1)	(2, 2)	(2,3)			
3	(3, 0)	(3, 1)	(3, 2)	(3, 3)			

**Theorem:** The rational numbers are countable

$$\mathbb{Q} = \{ a/b | a, b \in \mathbb{N}, b \neq 0 \}$$

**Hint:** Go through all possible numbers that the numerator and denominator can add up to

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- numerator and denominator add up to 3
  - $ightharpoonup 2 \mapsto \frac{2}{1}$

$$\mathbb{Q} = \{ \sqrt[a]{b} | \, a,b \in \mathbb{N}, b \neq 0 \}$$

- $ightharpoonup 0 \mapsto 0$
- numerator and denominator add up to 2
  - $ightharpoonup 1 \mapsto 1/1$
- numerator and denominator add up to 3
  - ightharpoonup 2  $\mapsto$  2/1
  - $\rightarrow$  3  $\mapsto$  1/2

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- numerator and denominator add up to 4

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	1	2	3	4	
1	1/1	1/2	1/3	1/4	
2	2/1	1	2/3	2/4	Ä
3	3/1	3/2	1	3/4	:
4	4/1	2	4/3	1	

Skip over redundant combinations

**Proposition:** The set of all possible <u>finite</u> binary strings is countable

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#### Uncountal

- AFSC biject
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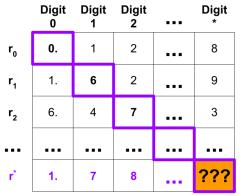
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### Diagonalization of $\mathbb R$



The assumption that  $\mathbb{R}$  is contable allows us to list out all of the real numbers (and then construct a paradoxical number)

Construct r\* by modifying the diagonals digits until we reach a contradiction

# Uncountability of infinite binary strings

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- ► Hint: proceed by contradiction
- construct a binary string that causes problems

► AFSOC the inifinite binary strings are countable. Then there is a bijection with N

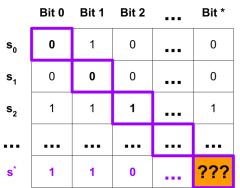
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# Diagonalization of Infinite Binary Strings



The assumption that the infinite binary strings are countable allows us to list out all of the infinite binary strings (and then construct a paradoxical binary string)

Construct s\* by modifying the diagonals bits until we reach a contradiction

► **Proposition:** The set of formal languages on any finite alphabet is uncountable

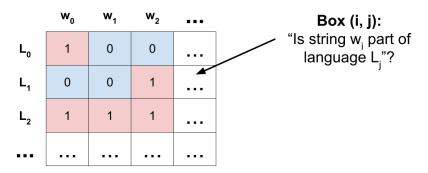
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- Let  $\Sigma^* = \{w_1, w_2, ...\}$  be the set of all possible stirngs
- ▶ Represent a language  $L_j \subseteq \Sigma^*$  using it's characteristic binary string  $S_j$

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- Let  $\Sigma^* = \{w_1, w_2, \dots\}$  be the set of all possible stirngs
- ▶ Represent a language  $L_j \subseteq \Sigma^*$  using it's characteristic binary string  $S_j$ 
  - ▶ The i-th digit of  $S_i$  s 1 if  $w_i \in L_i$  and 0 otherwise

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- ► The set of formal languages has a bijection with an uncountable set; thus it must be uncountable



We can represent every language using an infinite binary string

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- So there cannot possibly be a Turing machine to recognize every language