Arjun Chandrasekhar

A model of computation that uses a constant amount of memory to recognize a language.

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- Several states
 - One start state

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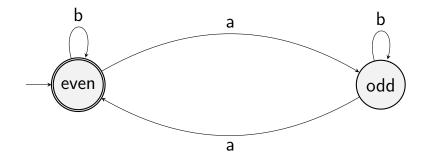
- One start state
- Some states are labelled as accept states

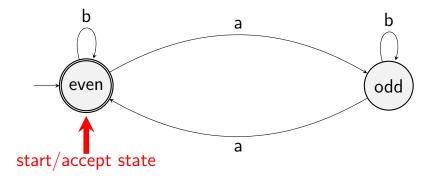
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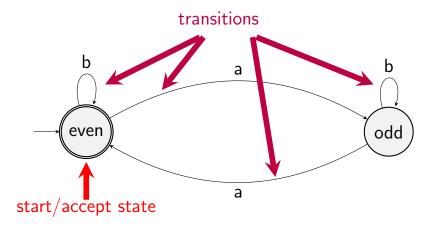
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- Some states are labelled as accept states
- Transitions (arrows) between states

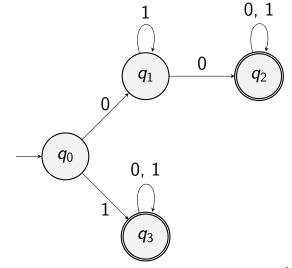
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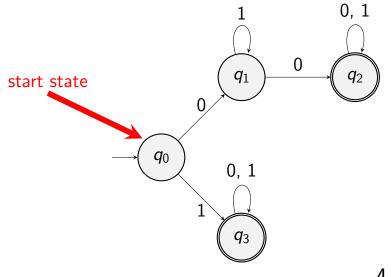
- One start state
- Some states are labelled as accept states
- Transitions (arrows) between states
 - Each transition is labelled by a character from the alphabet



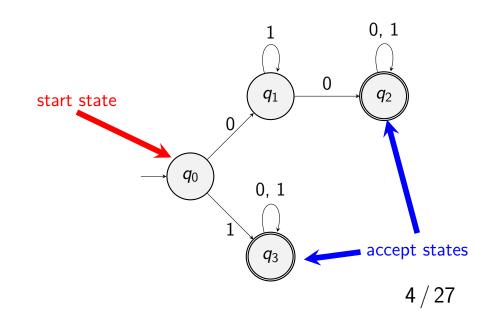


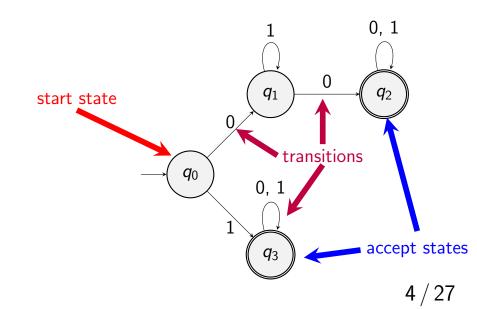


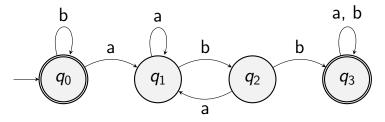




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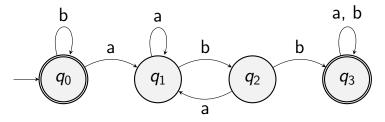




What is the start state of this DFA?

A. q_0 **C.** q_2

B. *q*₁ **D.** *q*₃

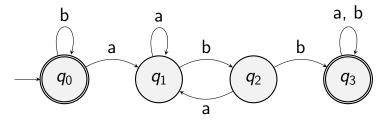


What is the start state of this DFA?

Α.	$q_0 \checkmark$	C.	q_2
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B. *q*₁ **D.** *q*₃



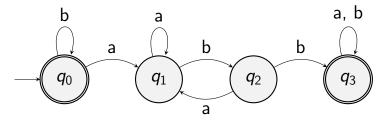


What are the accept states of this DFA?

A. q_0 **C.** q_2

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What are the accept states of this DFA?

A. $q_0 \checkmark$	C . <i>q</i> ₂
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B. q_1 **D.** $q_3 \checkmark$

Start in the start state

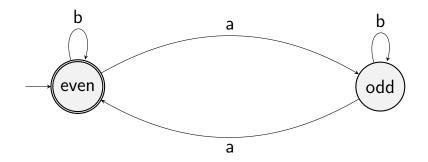
- Start in the start state
- Scan the symbols in the input one by one

- Start in the start state
- Scan the symbols in the input one by one
- For each symbol σ scaned:

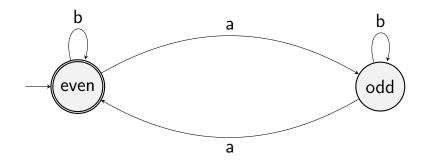
- Start in the start state
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- Start in the start state
- Scan the symbols in the input one by one
- For each symbol σ scaned:
 - \blacktriangleright Go to the next state by following the arrow with the label σ
- After scanning all of the input, if the DFA is in an accept state, the input is accepted. Otherwise the input is rejected

What does this DFA do on input aaaa?

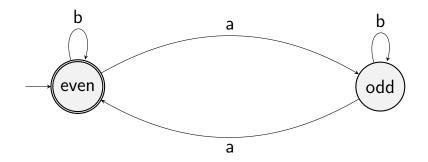


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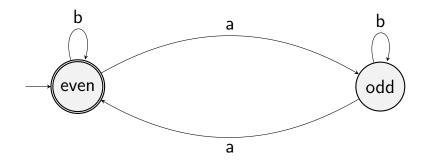


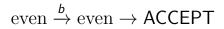
 $\operatorname{even} \xrightarrow{a} \operatorname{odd} \xrightarrow{a} \operatorname{even} \xrightarrow{a} \operatorname{odd} \xrightarrow{a} \operatorname{even} \rightarrow \mathsf{ACCEPT}$

What does this DFA do on input b?

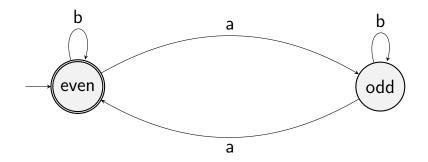


What does this DFA do on input b?

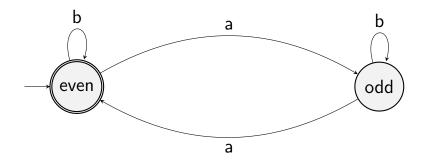




What does this DFA do on input abb?

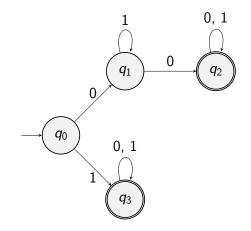


What does this DFA do on input abb?

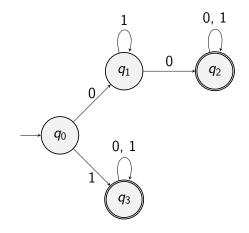


even \xrightarrow{a} odd \xrightarrow{b} odd \xrightarrow{b} odd \rightarrow REJECT

Deterministic Finite Automata What does this DFA do on input 0110?

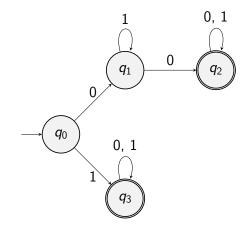


Deterministic Finite Automata What does this DFA do on input 0110?



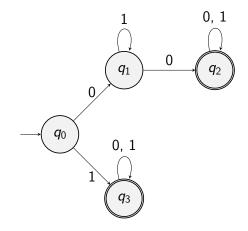
 $q_0 \xrightarrow{0} q_1 \xrightarrow{1} q_1 \xrightarrow{1} q_1 \xrightarrow{0} q_2
ightarrow \mathsf{ACCEPT} \ 11/27$

Deterministic Finite Automata What does this DFA do on input 1011?



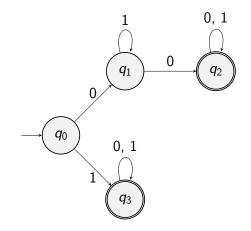
 $12 \, / \, 27$

Deterministic Finite Automata What does this DFA do on input 1011?

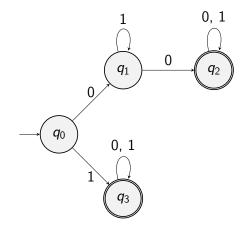


 $q_0 \xrightarrow{1} q_3 \xrightarrow{0} q_3 \xrightarrow{1} q_3 \xrightarrow{1} q_3 \rightarrow \mathsf{ACCEPT}$ 12/27

Deterministic Finite Automata What does this DFA do on input 0111?

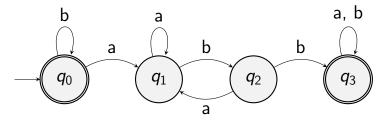


Deterministic Finite Automata What does this DFA do on input 0111?



 $q_0 \xrightarrow{0} q_1 \xrightarrow{1} q_1 \xrightarrow{1} q_1 \xrightarrow{1} q_1 \xrightarrow{1} q_1
ightarrow \mathsf{REJECT}$ 13 / 27

Deterministic Finite Automata



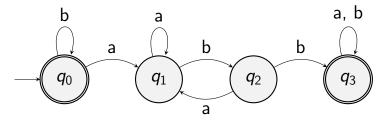
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A. ϵ C. bbbba

B. *abbbb*

D.
$$b^{100} = \underbrace{b \dots b}_{100}$$

Deterministic Finite Automata



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A. $\epsilon \checkmark$ C. bbbba

B. abbbb \checkmark **D.** $b^{100} = \underbrace{b \dots b}_{100} \checkmark$

$15 \, / \, 27$

Def: A **Deterministic Finite Automata (DFA)** is a 5-tuple $(Q, \Sigma, \delta, q_s, F)$

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- q_s the starting state
- F the set of accepting states

$16 \, / \, 27$

We say a DFA $D = (Q, \Sigma, q_s, \delta, F)$ accepts a string $w = \sigma_1 \sigma_2 \dots \sigma_n$ if there exists a sequence of states q_0, q_1, \dots, q_n such that:

16

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 q_i = δ(q_{i-1}, σ_i) for all i (all transitions are valid)

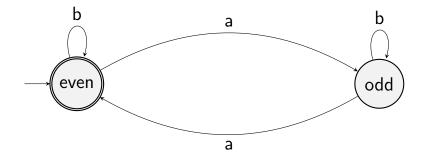
16

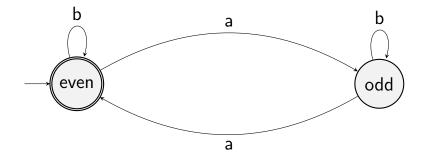
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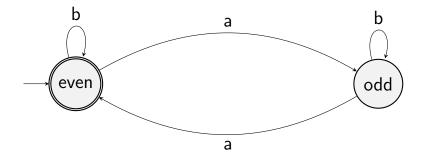
16

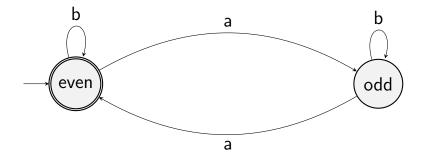
• $q_n \in F$ (finish in an accept state)





 $\blacktriangleright Q = \{even, odd\}$

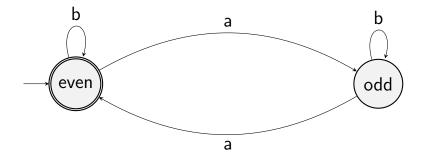




17

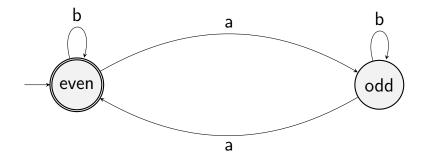
27

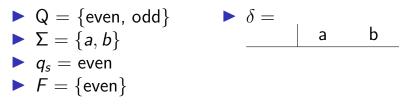
Q = {even, odd}
 Σ = {a, b}
 q_s = even

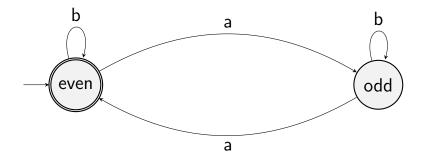


17

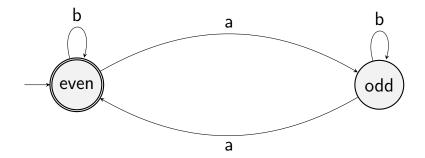
Q = {even, odd}
 Σ = {a, b}
 q_s = even
 F = {even}



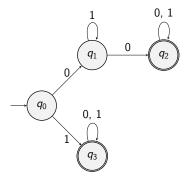


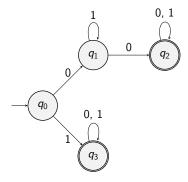


 $\begin{array}{l} \mathsf{Q} = \{ \mathsf{even}, \mathsf{odd} \} \\ \mathsf{\Sigma} = \{ a, b \} \\ \mathsf{q}_s = \mathsf{even} \\ \mathsf{F} = \{ \mathsf{even} \} \end{array}$ $\begin{array}{l} \mathsf{\delta} = \\ \underline{ \ } \\ \mathsf{even} \ \mathsf{odd} \ \mathsf{even} \\ \mathsf{even} \end{array}$

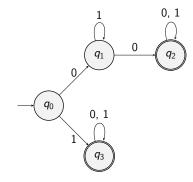


\blacktriangleright Q = {even, odd}	$\blacktriangleright \delta =$		
$\blacktriangleright \Sigma = \{a, b\}$		а	b
$\blacktriangleright q_s = even$	even	odd	even
$\blacktriangleright F = \{\text{even}\}\$	odd	even	odd



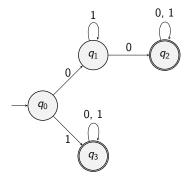


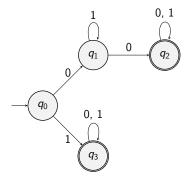
►
$$Q = \{q_0, q_1, q_2, q_3\}$$



•
$$Q = \{q_0, q_1, q_2, q_3\}$$

• $\Sigma = \{0, 1\}$



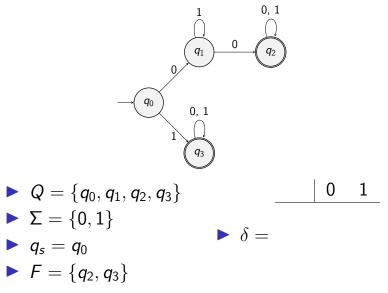


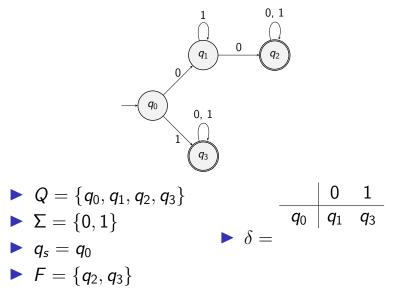
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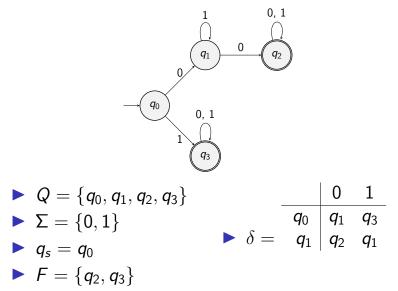
$$\Sigma = \{0, 1\}$$

$$q_s = q_0$$

$$F = \{q_2, q_3\}$$







$$D_{q_{1}} = \{q_{0}, q_{1}, q_{2}, q_{3}\}$$

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Write the formal definition

$$\begin{array}{c} & 1 & 0 & 1 \\ \hline q_1 & 0 & q_2 \\ \hline q_0 & 0, 1 \\ \hline q_3 \end{array}$$

$$\begin{array}{c} Q = \{q_0, q_1, q_2, q_3\} \\ \blacktriangleright \Sigma = \{0, 1\} \\ \blacktriangleright q_s = q_0 \\ \vdash F = \{q_2, q_3\} \end{array}$$

$$\begin{array}{c} 0 & 1 \\ \hline q_0 & q_1 & q_3 \\ \hline q_2 & q_2 & q_1 \\ \hline q_2 & q_2 & q_2 \\ \hline q_3 & q_3 & q_3 \end{array}$$

$19 \, / \, 27$

Let D be a DFA

$19 \, / \, 27$

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19

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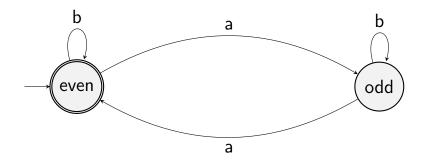
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 - Do we ever have to worry about D looping?

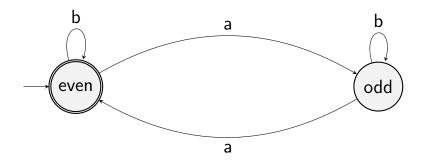
$19 \, / \, 27$

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 - If $w \in L$, then D accepts w
 - ▶ If $w \notin L$, then D rejects w
 - Do we ever have to worry about D looping?
- The language of D, denoted L(D) is the (unique) language that D recognizes - that is, the set of all strings that D accepts

What the language of this DFA?

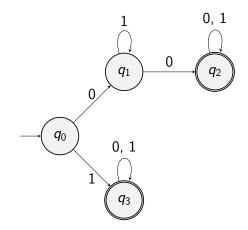


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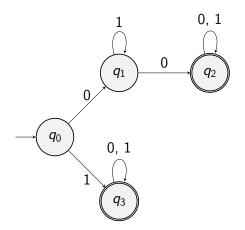


 $L(D) = \{w \mid w \text{ has an even number of a's}\}$

The language of a DFA What the language of this DFA?

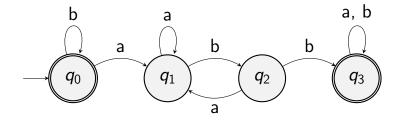


The language of a DFA What the language of this DFA?

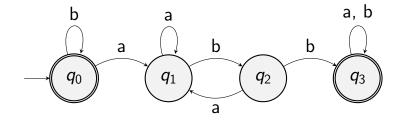


 $L(D) = \{w \mid w \text{ starts with 1 or has at least two 0's}$ 21 / 27

What the language of this DFA?



What the language of this DFA?



 $L(Q) = \{w \mid w \text{ contains only b's or contains abb}\}$



Some comments:

A transition must be defined for every symbol in the alphabet, for every state in the DFA.

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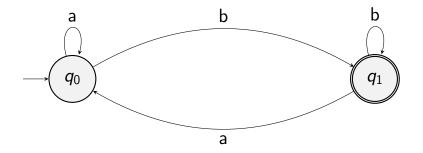
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- There can be more than one accepting state.
- The set of languages that can be recognized by some DFA is called the regular languages.
 - A language L is regular if and only if some DFA D recognizes L

Let $\Sigma = \{a, b\}$. Let's show that the following language is regular:

 $L = \{w \mid w \text{ ends with } b\}$

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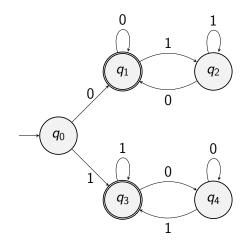


Let $\Sigma=\{0,1\}.$ Let's show that the following language is regular:

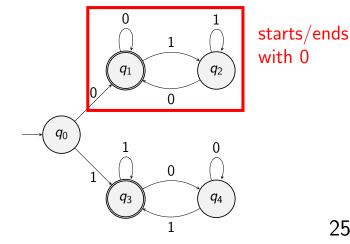
 $\mathsf{L} = \{\mathsf{w} \mid \mathsf{w} \text{ starts and ends with the same symbol}\}$

25

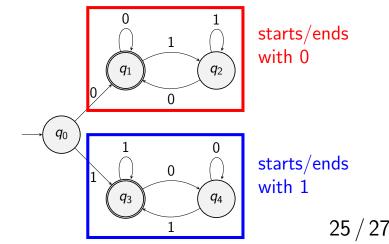
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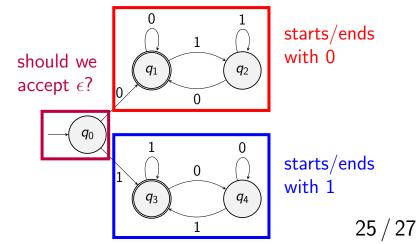
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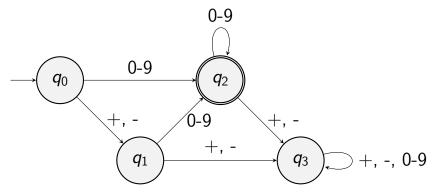


Design a DFA Let $\Sigma = \{+, -, 0, 1, \dots, 9\}$. Let's show that the following language is regular:

 $L = \{w \mid w \text{ is a valid integer literal}\}$

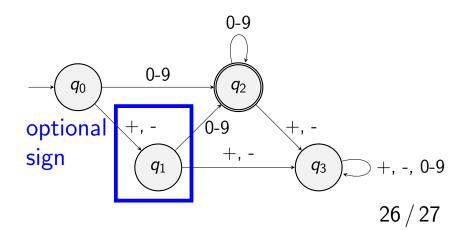
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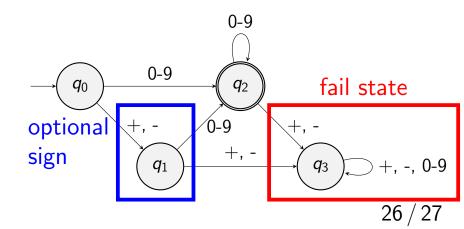
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